

Synchronous games and game algebras

Let $\mathcal{G} = (X, X, A, A, \lambda)$ be a **synchronous game**, i.e.

$$\lambda(x, x, a, b) = 0 \text{ if } a \neq b,$$

(e.g. $\text{Hom}(G, H)$) then winning strategies are characterised by traces on

$$\mathcal{A}(\lambda) = \mathcal{A}_{X,A}/J(\lambda) \text{ (the game algebra)}$$

with $J(\lambda) = \langle e_{x,a}e_{y,b} : \lambda(x, y, a, b) = 0 \rangle$.

We have (Helton-Meyer-Paulsen-Satriano, Kim-Paulsen-Schafhauser):

- \mathcal{G} has a **perfect C_{qc} strategy** iff there exists a non-zero tracial C^* -algebra (\mathcal{A}, τ) and a $*$ -homomorphism $\pi : \mathcal{A}(\lambda) \rightarrow \mathcal{A}$. Then

$$p(a, b|x, y) = \tau(e_{x,a}e_{y,b}),$$

- **C_{qa} strategy** iff we can take $\mathcal{A} = \mathcal{R}^U$;
- **C_q strategy** iff we can take \mathcal{A} finite-dimensional;
- **C_{loc} strategy** iff we can take \mathcal{A} abelian.

Homomorphism and isomorphism games

For $t \in \{loc, q, qa, qc\}$ and graphs G and H , write $G \rightarrow_t H$ if there exists a perfect C_t strategy for the homomorphism game $\text{Hom}(G, H)$. One can define quantum analogues of chromatic numbers and independence numbers $\chi(G)$ and $\alpha(G)$:

- The t -chromatic number of a graph G is

$$\chi_t(G) = \min\{d : G \rightarrow_t K_d\}$$

- The t -independence number of G is

$$\alpha_t(G) = \max\{d : K_d \rightarrow_t \bar{G}\}.$$

- $\chi(G) = \chi_{loc}(G) \geq \chi_q(G) \geq \chi_{qa}(G) \geq \chi_{qc}(G),$
 $\alpha(G) = \alpha_{loc}(G) \leq \alpha_q(G) \leq \alpha_{qa}(G) \leq \alpha_{qc}(G).$

$\mathcal{A}(\text{Hom}(G, H)), \mathcal{A}(\text{Iso}(G, H))$

$\text{Hom}(G, H)$ games are rich enough to witness differences between C_q , C_{qa} and C_{qc} :

Dykema-Paulsen-Prakash (2019) using *Kruglyak-Rabanovich-Samoilenko* (2003):

$$\exists G : \chi_q(G) > \chi_{qa}(G)$$

Mancinska-Roberson-Varvitsiotis (2020) using $\text{MIP}^* = \text{RE}$ (2020):

$$\exists G : \alpha_{qc}(G) > \alpha_{qa}(G)$$

$\text{Iso}(G, H)$ games are examples of bisynchronous games where

$$\lambda(x, x, a, b) = 0 \text{ and } \lambda(x, y, a, a) = 0 \text{ if } x \neq y, a \neq b.$$

Perfect strategies are captured by **quantum permutation group**: $O(S_X^+)$ is generated by $p_{a,x} = p_{a,x}^* = p_{a,x}^2$ s.t. $U = (p_{a,x})_{a,x \in X}$ is a magic unitary, i.e. $\sum_a p_{a,x} = \sum_y p_{b,y} = 1$, $x, b \in X$.

If A_G and A_H are the adjacency matrices of G and H resp. then

$$\mathcal{A}(\text{Iso}(G, H)) = \langle p_{a,x} : U = (p_{a,x})_{a,x} \text{ magic unitary with } (A_G \otimes 1)U = U(A_H \otimes 1) \rangle.$$

Combining the previous results one gets:

For $\text{Iso}(G, H)$ we have:

- $G \simeq_{qc} H$ iff there is a trace on $\mathcal{A}(\text{Iso}(G, H))$
- $G \simeq_q H$ iff there is a fin.-dim. repr. of $\mathcal{A}(\text{Iso}(G, H))$
- $G \simeq_{loc} H$ iff $G \simeq H$ iff there is a one-dim repr. of $\mathcal{A}(\text{Iso}(G, H))$.

$$G \simeq H \xRightarrow[\neq]{\text{[AMRSSV]}} G \simeq_q H \xRightarrow[\neq]{\text{[KPS]}} G \simeq_{qa} H \Rightarrow G \simeq_{qc} H$$

$$G \simeq_{qc} H \Leftrightarrow \begin{matrix} G \simeq_{C^*} H \\ * \text{Rep} \mathcal{A}(\text{Iso}(G, H)) \neq \emptyset \end{matrix} \Leftrightarrow \begin{matrix} G \simeq_{alg} H \\ \mathcal{A}(\text{Iso}(G, H)) \neq \emptyset \end{matrix} \quad (\text{[BCEHPSW]})$$

Going quantum-to-quantum

Motivating question I: Suppose the game has **quantum inputs/outputs**.

What kind of strategies can be used?

Motivating question II: Perhaps a suitable (**simpler and genuinely**) **quantum game** can disprove Tsirelson-Connes?

A classical input $(x, y) \rightsquigarrow$ the state $\epsilon_{x,x} \otimes \epsilon_{y,y} \in \mathcal{D}_X \otimes \mathcal{D}_Y$ (pos.-semidef. matrix of trace 1) here **matrix units**: $\epsilon_{x,y}, x, y \in X$ in M_X .

A correlation $p = \{(p(a, b|x, y))_{a,b} : x, y \in X\} \rightsquigarrow \mathcal{N}_p : \mathcal{D}_{X \times Y} \rightarrow \mathcal{D}_{A \times B}$,

$$\mathcal{N}_p(\epsilon_{x,x} \otimes \epsilon_{y,y}) = \sum_{a,b \in A} p(a, b|x, y) \epsilon_{a,a} \otimes \epsilon_{b,b}.$$

Note: p is no-signalling \Leftrightarrow

$$\mathrm{Tr}_A \mathcal{N}_p(\epsilon_{x,x} \otimes \epsilon_{y,y}) = \mathrm{Tr}_A \mathcal{N}_p(\epsilon_{x',x'} \otimes \epsilon_{y,y}) \quad \text{and} \quad \mathrm{Tr}_B \mathcal{N}_p(\epsilon_{x,x} \otimes \epsilon_{y,y}) = \mathrm{Tr}_B \mathcal{N}_p(\epsilon_{x,x} \otimes \epsilon_{y',y'})$$

Quantisation: (*Duan-Winter*) **Quantum channels** (completely positive trace preserving)

$\Gamma : M_{XY} \rightarrow M_{AB}$, satisfying no-signalling conditions.

No-signalling:

$$\mathrm{Tr}_A \Gamma(\omega_X \otimes \omega_Y) = \mathrm{Tr}_A \Gamma(\omega'_X \otimes \omega_Y) \quad \text{and} \quad \mathrm{Tr}_B \Gamma(\omega_X \otimes \omega_Y) = \mathrm{Tr}_B \Gamma(\omega_X \otimes \omega'_Y)$$

Classes of quantum no-signalling correlations (Todorov-T. (2020))

A family of classical POVM's:

$$\{(E_{x,a})_{a \in A} : x \in X\}$$

\rightsquigarrow

$$E = \sum_{x \in A} \sum_{a \in A} \epsilon_{x,x} \otimes \epsilon_{a,a} \otimes E_{x,a} \in M_{XA}(B(H))^+$$

A family of quantum POVM's:

Stochastic operator matrix

$$E = (E_{x,x',a,a'}) \in M_{XA}(B(H))^+ \text{ such that } \text{Tr}_A E = I \otimes I_X.$$

Classes of quantum no-signalling (QNS) correlations

\updownarrow
 Quantum channels $\Gamma : M_{XY} \rightarrow M_{AB}$ with specified Choi matrices
 $(\Gamma(\epsilon_{x,x'} \otimes \epsilon_{y,y'}))_{x,x',y,y'} \in M_{XYAB}$

Local QNS

Convex combinations of $\Phi \otimes \Psi$

Quantum QNS

$$(\langle E_{x,x',a,a'} \otimes F_{y,y',b,b'} \xi, \xi \rangle), \quad \xi \in H_A \otimes H_B$$

Quantum commuting QNS

$$(\langle E_{x,x',a,a'} F_{y,y',b,b'} \xi, \xi \rangle), \quad \xi \in H$$

$$\mathcal{Q}_{\text{loc}} \subset \mathcal{Q}_{\text{q}} \subset \mathcal{Q}_{\text{qa}} \subset \mathcal{Q}_{\text{qc}} \subset \mathcal{Q}_{\text{ns}}$$

Quantum non-local games

Classical games (X, Y, A, B, λ) :

Rule function $\lambda : X \times Y \times A \times B \rightarrow \{0, 1\} \rightsquigarrow (P_{(x,y)}, P_{\beta_{x,y}(\lambda)})$

$\beta_{x,y}(\lambda) = \{(a, b) : \lambda(x, y, a, b) = 1\}$,

P_α is a projection onto $\text{span}\{e_x \otimes e_y : (x, y) \in \alpha\}$

$\lambda \rightsquigarrow \varphi_\lambda : \text{Proj}(\mathcal{D}_{XY}) \rightarrow \text{Proj}(\mathcal{D}_{AB})$, V -preserving 0-preserving map.

Quantum games:

Replace (X, Y, A, B, λ) by $\text{Proj}(M_{XY}), \text{Proj}(M_{AB})$ and

$$\varphi : \text{Proj}(M_{XY}) \rightarrow \text{Proj}(M_{AB}),$$

V -preserving 0-preserving map.

Definition (Todorov-T, 2020)

A QNS correlation $\Gamma : M_{XY} \rightarrow M_{AB}$ is a perfect strategy for the quantum non-local game φ if $\langle \Gamma(P), \varphi(P)^\perp \rangle := \text{Tr}(\Gamma(P)\varphi(P)^\perp) = 0, \forall P \in \text{Proj}(M_{XY})$.

Questions and some references

- "Synchronous" and "bisynchronous" quantum games and associated algebras.

Brannan-Harris-Todorov-T, 2021,2022, Bochniak-Kaspzak-Sołtan, 2021.

- Quantum graph homomorphisms and isomorphisms and associated algebras.

Brannan-Chirvasitu-Eifler-Harris-Paulsen-Su-Wasilewski 2020,

Brannan-Ganesan-Harris 2020, Ganesan 2022, Brannan-Harris-Todorov-T, 2021,2022

- Value of quantum non-local games.

Cooney-Junge-Palazuelos-Pérez-García 2015, Crann-Levene-Todorov-T-Winter, in progress